

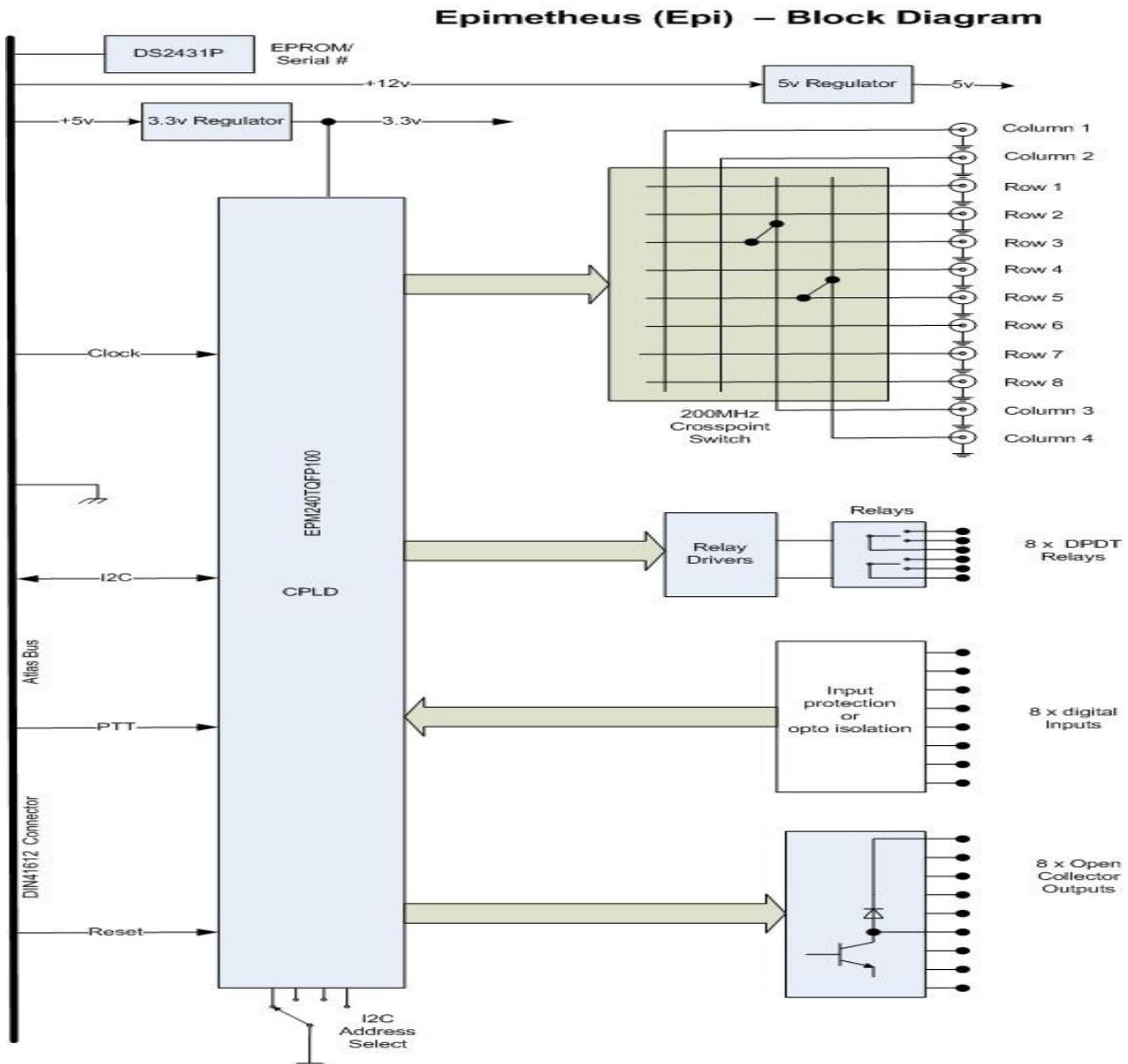
EPIMETHEUS FPGA/CPLD Specification



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1 EPIMETHEUS FPGA/CPLD Specification

This document describes the Verilog code which is the core of the Epimetheus (EPI) CPLD functionality. The board block diagram is shown below:



EPIMETHEUS Block Diagram

Epi V1.0 27 September 2006

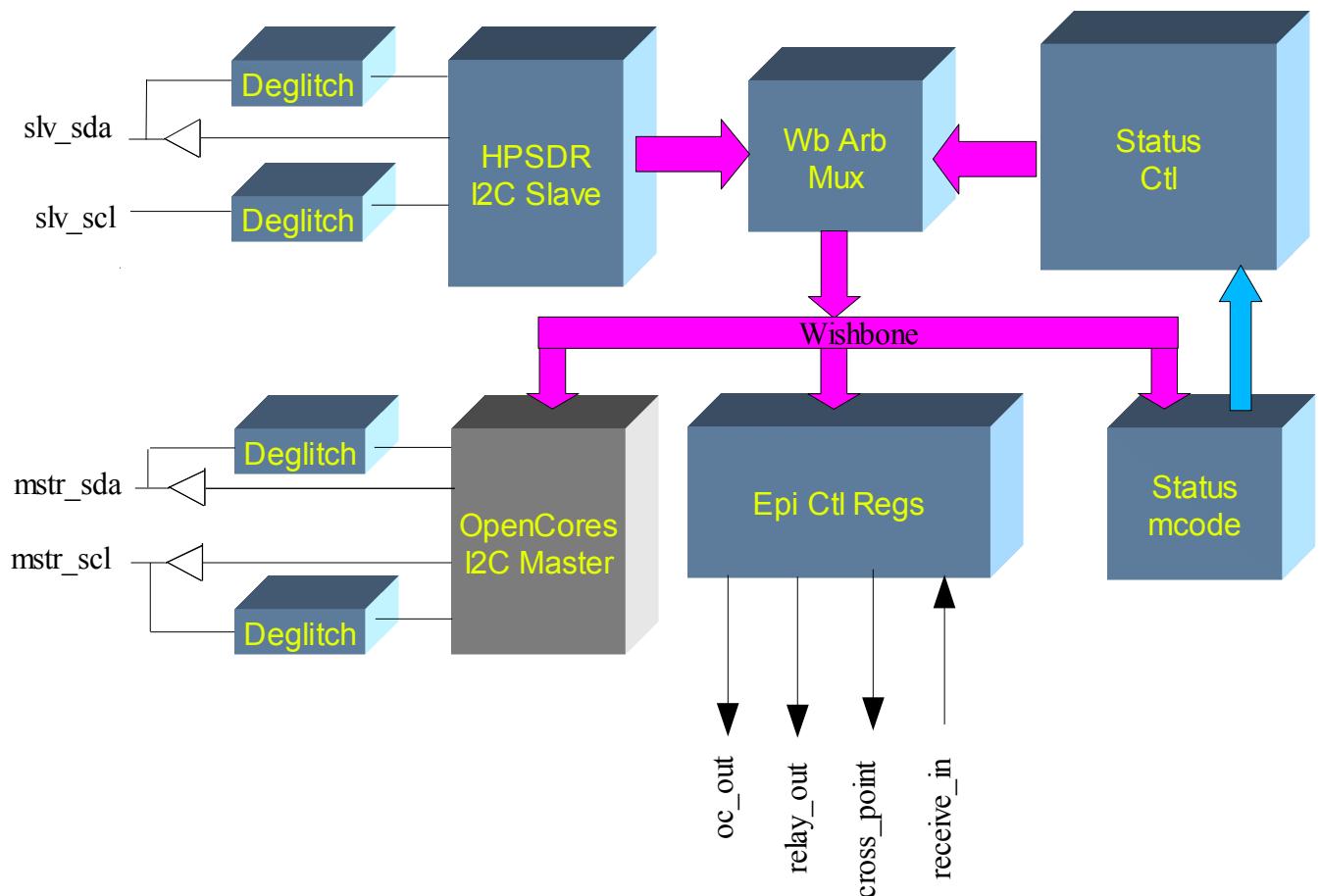
2 Functional Description

The main mission of the Epi CPLD is to provide programmed Outputs for the various facilities on EPI and to be capable of sending a message to the host processor on Ozy whenever a state change occurs on the 'receive_in' pins (like a button push).

The I2C bus is used both as the programming path AND the message path. There are two I2C ports on the device primarily to simplify internal tristate control. These should be tied together on a single I2C bus on the EPI board.

The Verilog code for Epimetheus implements an I2C slave which translates I2C activity into Wishbone bus cycles. These Wishbone transactions can target one of three slaves, the I2C Master (from OpenCores), the "Epi_ctl_regs" module, and the "Status_mcode" module. The "Status_ctl" module serves as another Wishbone master which causes messages over the I2C master every time a signal changes state on the "receive_in" lines. The "Wb_arb_mux" module serves as an arbiter between the two Wishbone bus masters and feeds the data bus to/from each of the three I2C slaves.

A block Diagram of the EPIMETHEUS code is illustrated below:



Epi Verilog Code Block Diagram

This document covers all of the modules colored in Blue. The “OpenCores I2C Master” has it's own documentation that can be found on the “www.opencores.org” web page.

2.1 *EPI Module description*

This module is the top level and incorporates several sub-modules which implement the CPLD functionality. The table below illustrates the module I/Os.

Signal Name	Width	Direction	Description
reset_n	1	In	Reset (Active low)
CLK	1	In	Clock input – should be 30-40Mhz
slv_sda	1	In/Out	I2C Slave Serial Data
slv_scl	1	In	I2C Slave Serial Clock
mstr_sda	1	In/Out	I2C Master Serial Data
mstr_scl	1	In/Out	I2C Master Serial Clock
cross_point_ctl	48	Out	Cross Point Control
relay_out	8	Out	Relay Control
receive_in	8	In	Received Inputs
oc_out	8	Out	Open Collector Out

2.2 *OpenCores I2C Master*

The OpenCores I2C master is a wishbone slave that is employed within this design to send messages to the system host (Ozy) when any bit within “receive_in” changes state. The initial conditions are set up by the I2C slave from Ozy, then the microcode inside “Status_mcode” that drives “Status_ctl” will send a copy of the “receive_in” bus anytime a bit changes state.

The table below illustrates the module I/O – please see the OpenCores I2C Master document for further design details.

Signal Name	Width	Direction	Description
wb_clk_i	1	In	Wishbone Clock
wb_RST_i	1	In	Wishbone Reset
arst_i	1	In	Asynchronous reset
wb_adr_i	3	In	Wishbone Address
wb_dat_i	8	In	Wishbone Data In
wb_dat_o	8	Out	Wishbone Data Out

wb_we_i	1	In	Wishbone Write
wb_stb_i	1	In	Wishbone Strobe/Select
wb_cyc_i	1	In	Wishbone Cycle
wb_ack_o	1	Out	Wishbone Ack
wb_inta_o	1	Out	Wishbone Interrupt Out
scl_pad_i	1	In	I2C Clock In
scl_pad_o	1	Out	I2C Clock Out
scl_padoen_o	1	Out	I2C Clock Pad Out Enable (active low)
sda_pad_i	1	In	I2C Data In
sda_pad_o	1	Out	I2C Data Out
sda_padoen_o	1	Out	I2C Data Pad Out Enable (active low)

Please see the register description section for register bit definitions.

The define ADDRESS sets the I2C Slave address.

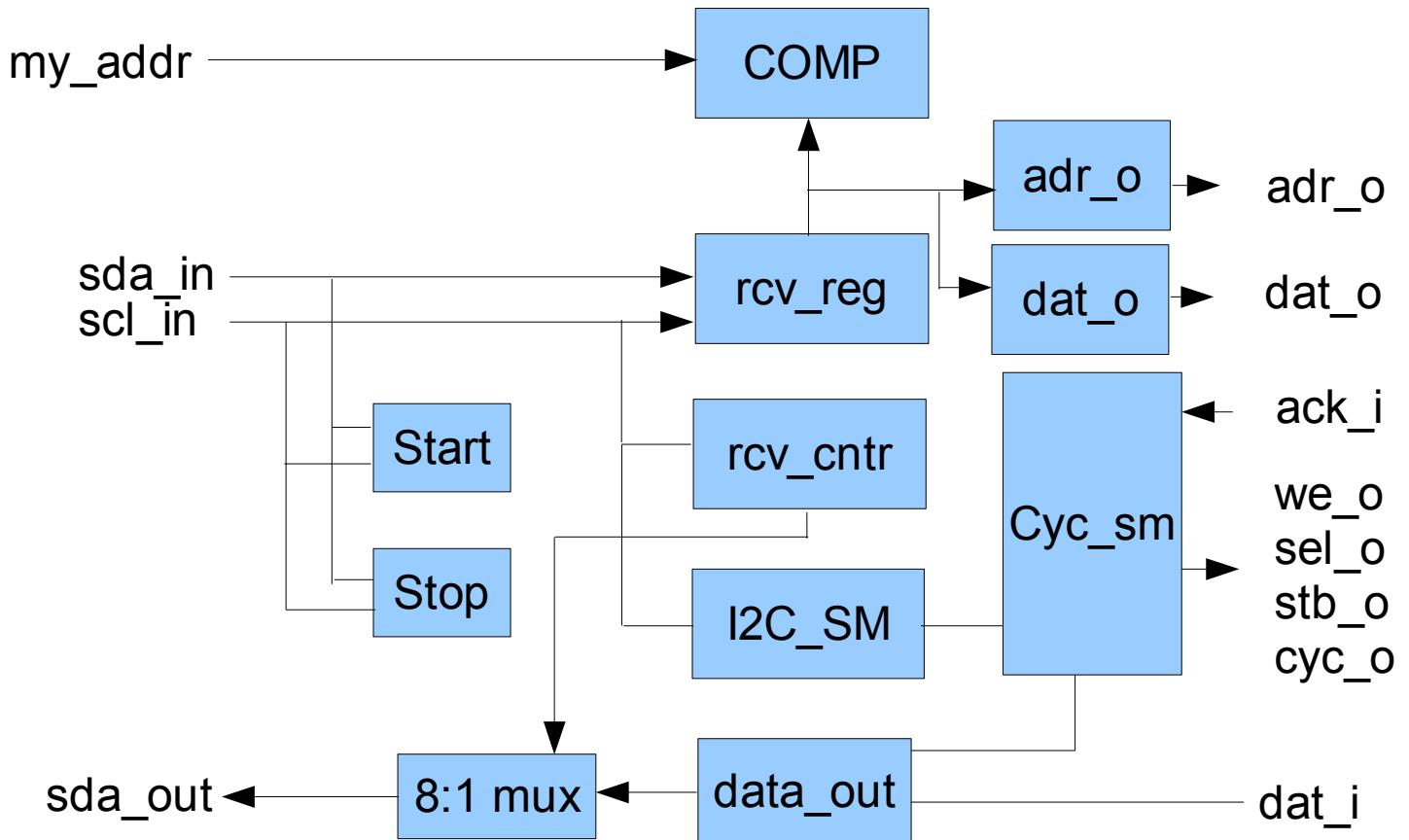
2.3 I2C_SLAVE (HPSDR)

This module implements an I2C slave which translates I2C transactions into Wishbone operations.

The module receives its I2C address from the top level module (the address which it responds on the I2C too). It is assumed that the I2C signals are de-glitched and that SDA occurs AFTER SCL to allow the STOP & START states to be determined reliably.

Signal Name	Width	Direction	Description
rst_n	1	In	reset (active low)
clk	1	In	Clock
my_addr	7	In	I2C Slave Address
adr_o	8	Out	Wishbone Address
dat_i	8	In	Wishbone Data In
dat_o	8	Out	Wishbone Data Out
we_o	1	Out	Wishbone Write
stb_o	1	Out	Wishbone Strobe
sel_o	1	Out	Wishbone Select (redundant with strobe)
cyc_o	1	Out	Wishbone Cycle
ack_i	1	In	Wishbone Ack
debug	12	Out	Debug bus (not required)

scl_in	1	In	I2C Clock In
sda_in	1	In	I2C Data In
sda_out	1	Out	I2C Data Out
sda_oe	1	Out	I2C Data Pad Out Enable (active high)



I2C_SLAVE Block Diagram

The “START” detector detects the transition of SDA from high to low while the clock is high.

The “STOP” detector sets when the transition of SDA from low to high occurs while the clock is high.

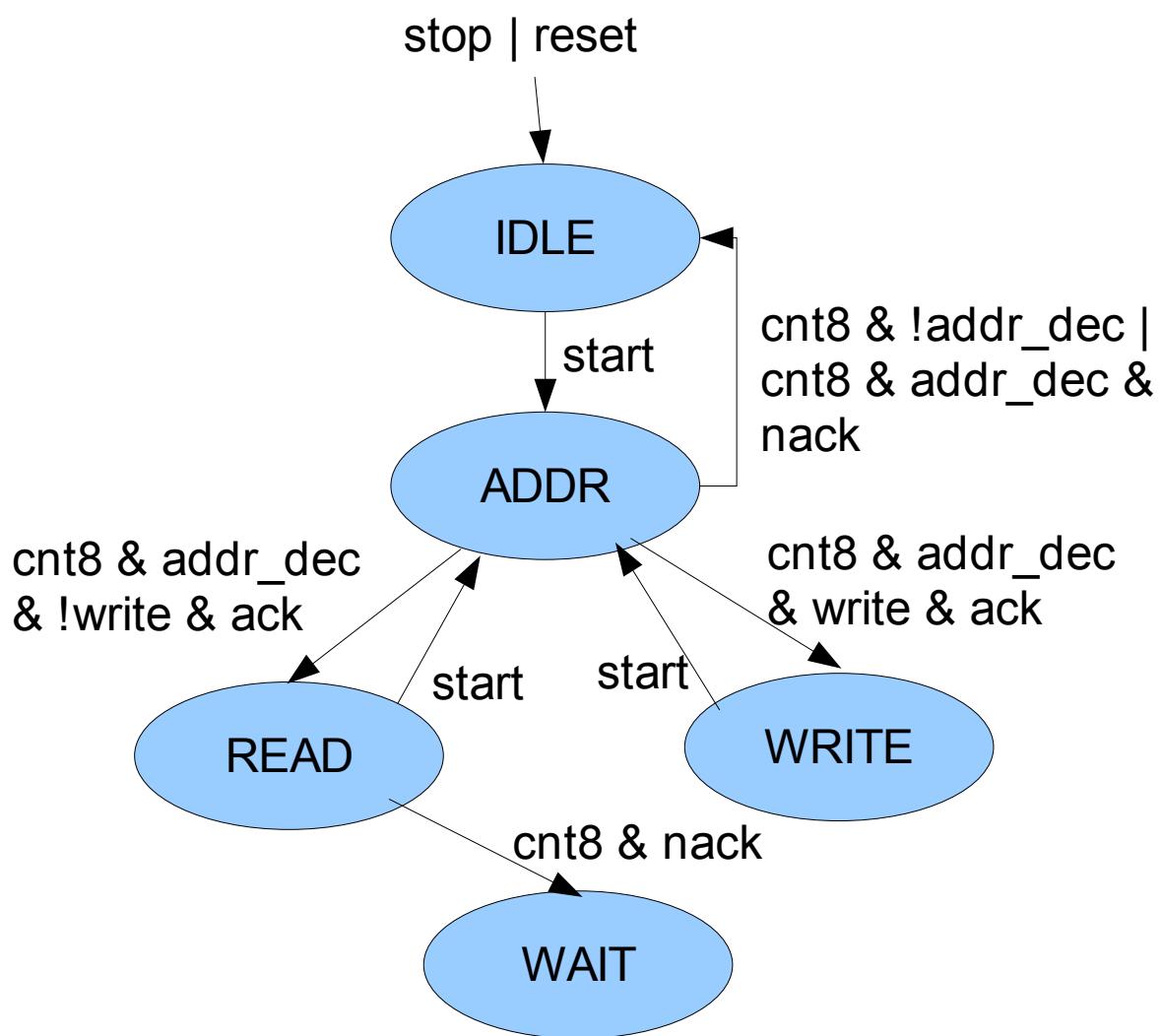
The “RCV_REG” shifts in the received data when the SCL transitions from low to high. This allows it to capture the data transmitted over the I2C bus.

The “RCV_CNTR” counts the number of bits that have been shifted into “RCV_REG.”

The “COMP” function watches the received data, and when the I2C_SM is in the SM_ADDR state and the data compares to the “my_addr” bus, it indicates an that device has been selected.

The “Data_out” module captures the data from the Wishbone dat_i bus and stores it to be sent out over the I2C SDA_OUT signal through the “8:1 Mux.”

The I2C_SM follows the state diagram below:



I2C SM State Diagram

The state machine goes asynchronously to IDLE when STOP occurs. This is how it exits the READ, WAIT, and WRITE states nominally during burst operations.

The function of the state machine is to interpret multiple I2C sequences, separating them into Address phase which selects the target I2C device followed by either Write or Read phases.

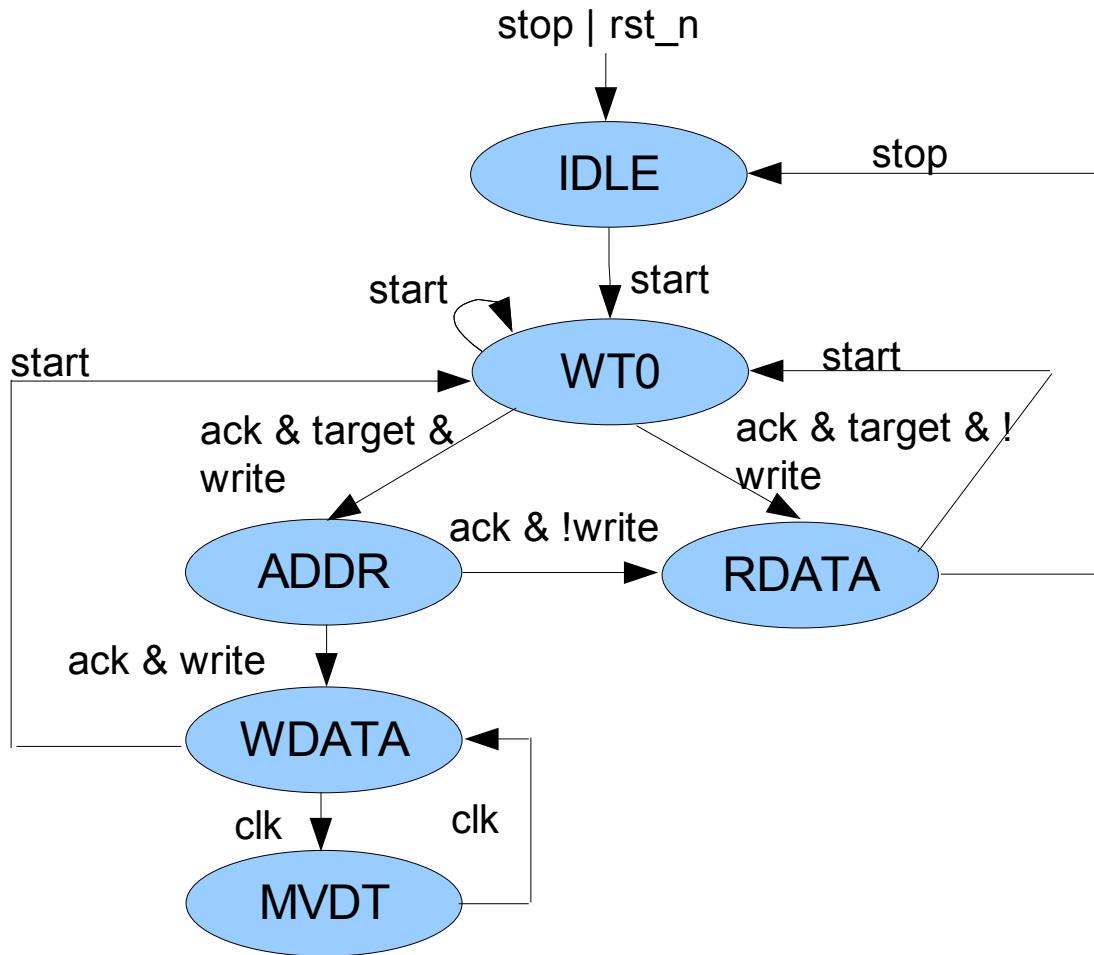
Start	I2C Addr+RW	a c k	Wishbone Address	a c k	Wishbone Write Data	a c k	Stop
IDLE	ADDR		WRITE		WRITE		IDLE

I2C State Machine Sequence to Write Data to the Wishbone bus

Start	I2C Addr+RW	a c k	Wishbone Address	a c k	Start	I2C Addr+RW	a c k	Wishbone Read Data	a c k	Stop
IDLE	ADDR		WRITE		ADDR			READ		

I2C State Machine Sequence to Read Data via the Wishbone bus

The “CYC_SM” is used to cause read and write transactions on the Wishbone bus. The CYC_SM interprets the different points within the I2C sequences and takes appropriate action to move data to or from the I2C resources across the Wishbone bus.



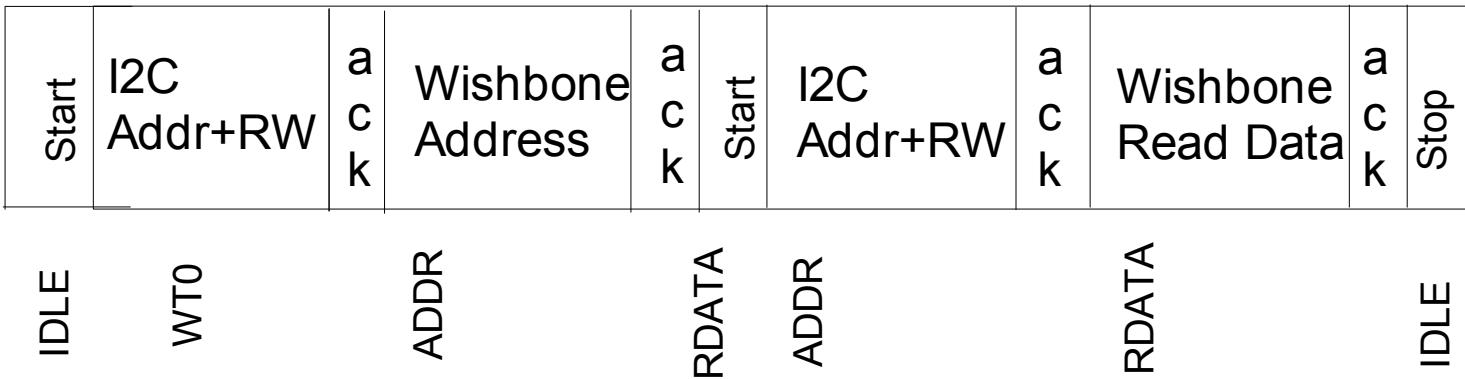
CYC_SM State Machine

Note that ALL of the signal state points mentioned above are synchronized copies of the state signals that exist in the I2C clock domain space. This is necessary to remove any problems with metastability.

Start	I2C Addr+RW	a c k	Wishbone Address	a c k	Wishbone Write Data	a c k	Stop
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CYC_SM state sequence following the I2C Write Transfer



CYC_SM State sequence following the I2C Read Transfer

The “ADR_O” register captures data upon the state machine moving from ADDR to WDATA. “ADR_O” will increment on succeeding occurrences of ACK while in the WDATA state or RDATA state. This supports burst operations for both read and write sequences.

The “DAT_O” register captures data upon the state machine moving when the state machine is in the WDATA state and ACK occurs.

Wishbone cycles occur when in the WDATA state and ACK occurs, in the RDATA state and the transfer count = 8, or during the WT0 period to allow pre-fetch of read data.

2.4 EPI_CONTROL_REGS Module

The Epi Control Regs module implements two basic functions. These are the Read/Write path to the programmable I/O pins and the detection of state change on the “received_in” pins. The module is a Wishbone slave and is nominally only accessible from the I2C Slave master.

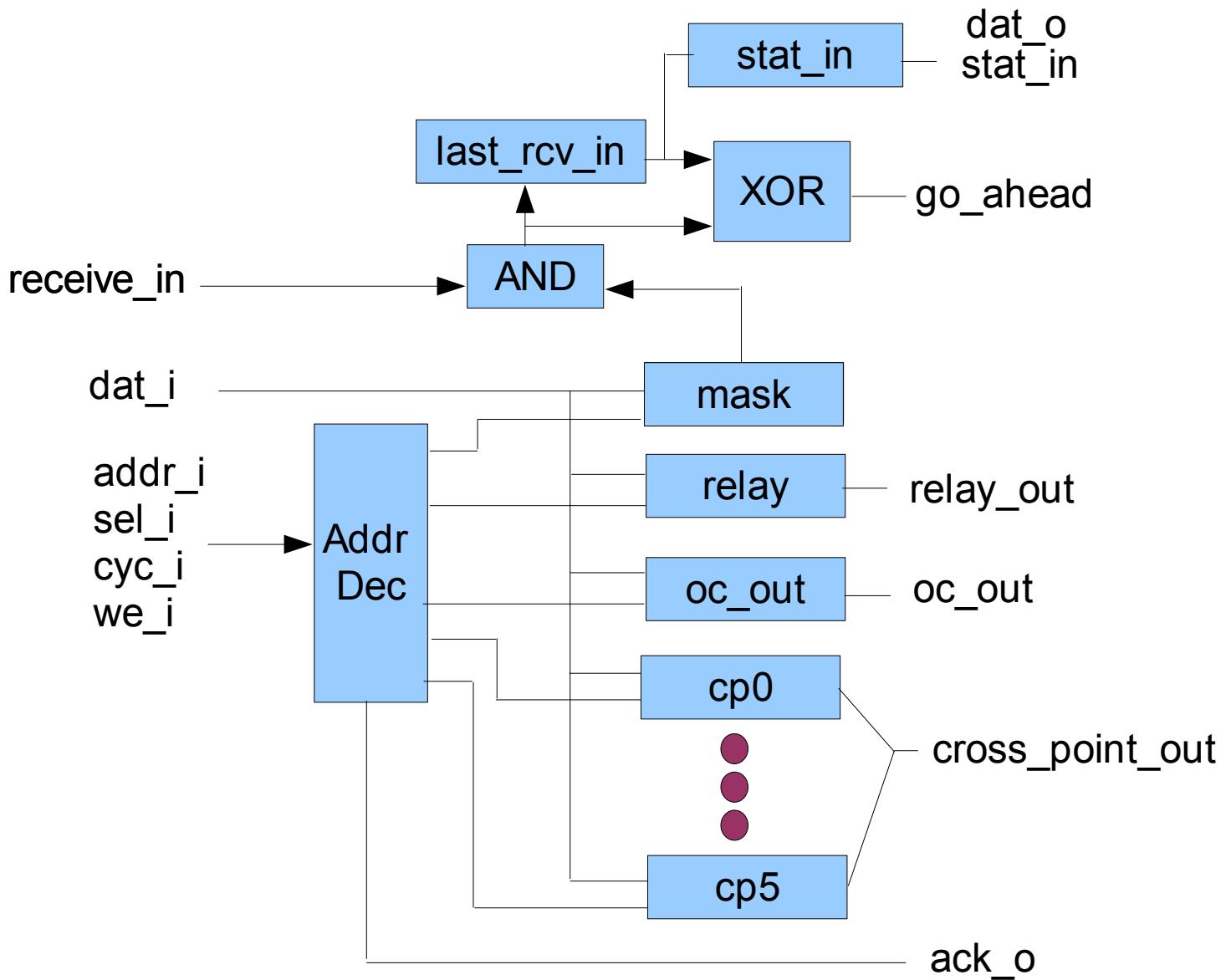
Signal Name	Width	Direction	Description
rst_n	1	In	reset (active low)
clk	1	In	Clock
addr_o	8	Out	Wishbone Address
dat_i	8	In	Wishbone Data In
dat_o	8	Out	Wishbone Data Out
we_i	1	Out	Wishbone Write
sel_i	1	Out	Wishbone Select (redundant with strobe)
cyc_i	1	Out	Wishbone Cycle
ack_o	1	In	Wishbone Ack
go_ahead	1	Out	Start status_ctl SM – a single clock pulse
cross_point_ctl	48	Out	Cross Point programmable control

relay_out	8	Out	Relay programmable control
oc_out	8	Out	Open Collector programmable Control
stat_in	8	Out	Masked received_in
receive_in	8	In	Received external bus

Write Addressable registers

Address	Reg Name	Description
8'h40	MASK	Masks the received_in bus when bit is 0.
8'h41	RELAY	Controls “relay_out[7:0]”
8'h42	OC_OUT	Controls “oc_out[7:0]”
8'h43	CP0	Controls “cross_point_out[7:0]”
8'h44	CP1	Controls “cross_point_out[15:8]”
8'h45	CP2	Controls “cross_point_out[23:16]”
8'h46	CP3	Controls “cross_point_out[31:24]”
8'h47	CP4	Controls “cross_point_out[39:32]”
8'h48	CP5	Controls “cross_point_out[47:40]”

Note all read accesses receive the STAT_IN bus which is a time delayed version of the XOR output.



EPI_CTL_REGS block diagram

There are 9 separate registers that can be written to, while a read to any address returns the “STAT_IN” bus.

The “receive_in” bus is ANDed with the value in the “MASK” register and delayed by one clock. The result of the masked value is compared to the delayed value to detect a change in state. Any change in state will cause the “go_ahead” pulse to occur.

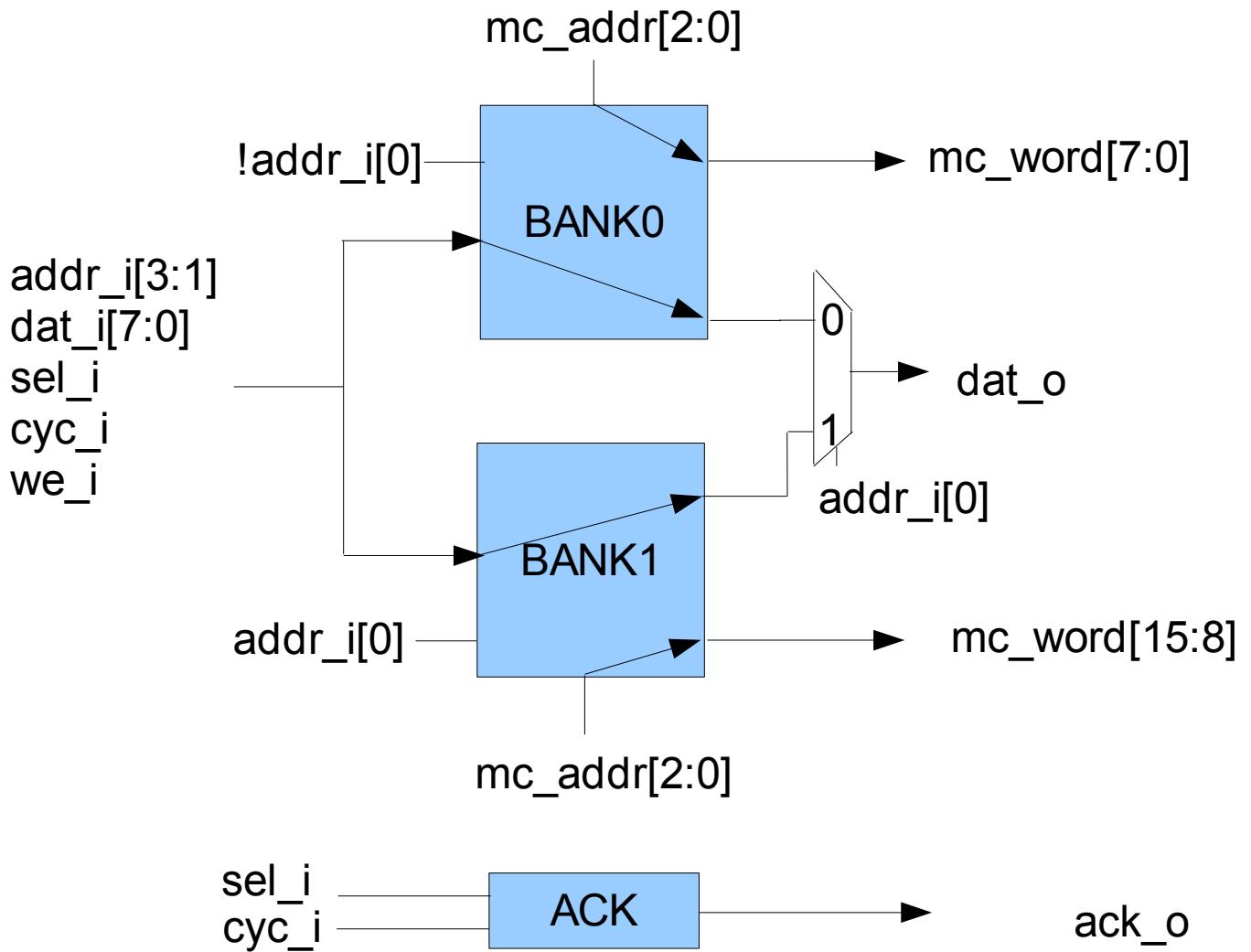
2.5 STATUS_MCODE

This module implements a byte addressable 16X8 memory array that is attached as a Wishbone slave to the internal Wishbone bus. Bank 0 receives even addresses, while Bank1 receives odd addresses. A second read port feeds the Microcode word over to the “STATUS_CTL” module.

Signal Name	Width	Direction	Description
rst_n	1	In	reset (active low)
clk	1	In	Clock
addr_i	8	Out	Wishbone Address
dat_i	8	In	Wishbone Data In
dat_o	8	Out	Wishbone Data Out
we_i	1	Out	Wishbone Write
sel_i	1	Out	Wishbone Select (redundant with strobe)
cyc_i	1	Out	Wishbone Cycle
ack_o	1	In	Wishbone Ack
mc_addr	3	In	Microcode Address
mc_word	16	Out	Microcode Data Word

Read/Write address descripton

Address	Reg Name	Description
8'h00	Bank0 byte 0	Microcode Word 0 – lower byte
8'h01	Bank1 byte 0	Microcode Word 0 – upper byte
8'h02	Bank0 byte 1	Microcode Word 1 – lower byte
8'h03	Bank1 byte 1	Microcode Word 1 – upper byte
8'h04	Bank0 byte 2	Microcode Word 2 – lower byte
8'h05	Bank1 byte 2	Microcode Word 2 – upper byte
8'h06	Bank0 byte 3	Microcode Word 3 – lower byte
8'h07	Bank1 byte 3	Microcode Word 3 – upper byte
8'h08	Bank0 byte 4	Microcode Word 4 – lower byte
8'h09	Bank1 byte 4	Microcode Word 4 – upper byte
8'h0A	Bank0 byte 5	Microcode Word 5 – lower byte
8'h0B	Bank1 byte 5	Microcode Word 5 – upper byte
8'h0C	Bank0 byte 6	Microcode Word 6 – lower byte
8'h0D	Bank1 byte 6	Microcode Word 6 – upper byte
8'h0E	Bank0 byte 7	Microcode Word 7 – lower byte
8'h0F	Bank1 byte 7	Microcode Word 7 – upper byte



STATUS Microcode store block diagram

The Wishbone bus organizes the memory into two banks, each 8 bits wide that take on even an odd addresses respectively on the Wishbone bus. The second read path is addressed by $mc_addr[2:0]$ and supplies the “ mc_word ” bus.

2.6 STATUS_CTL

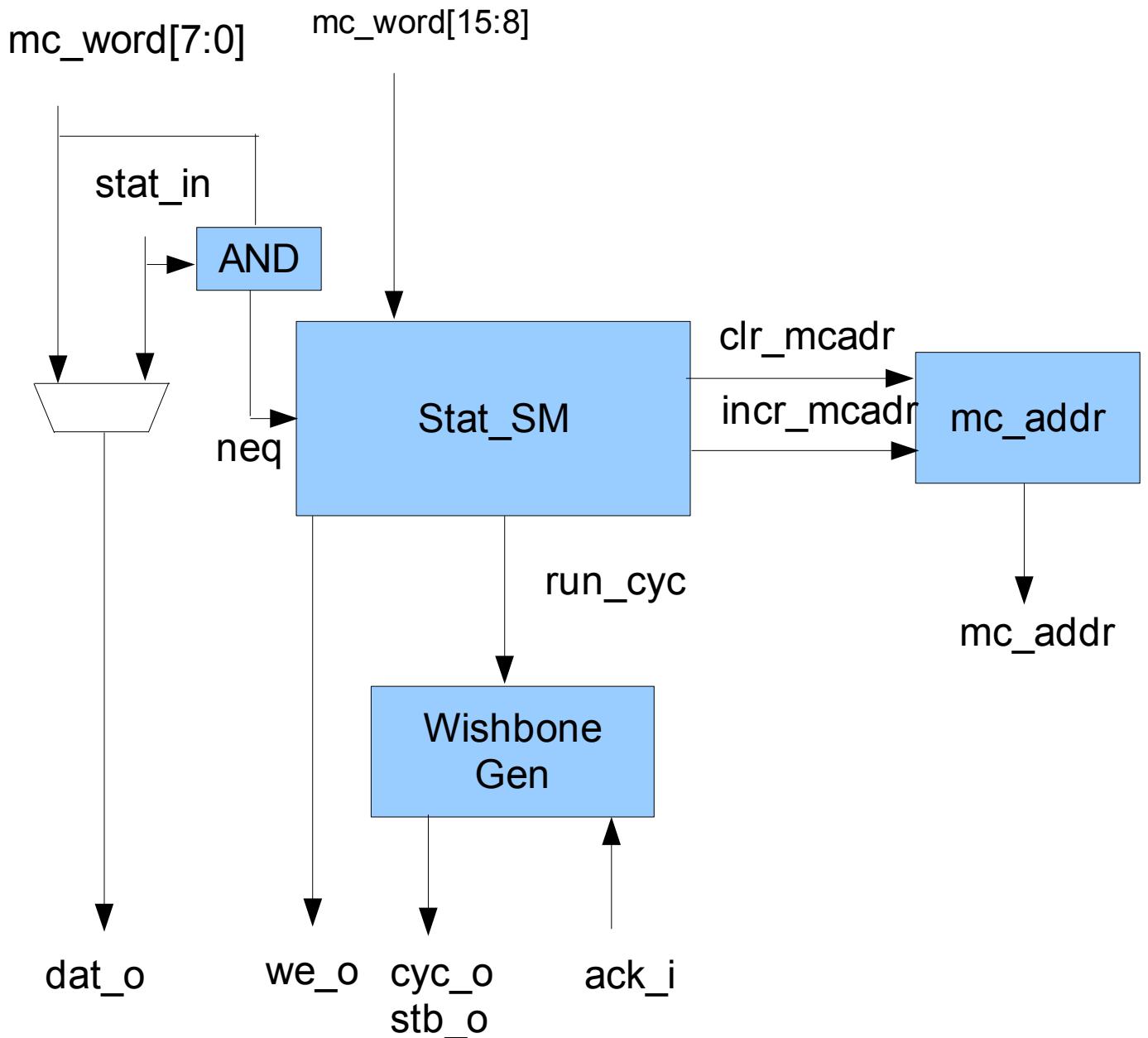
The “STATUS_CTL” module is serves as a Wishbone master which causes I2C messages to be sent via the I2C Master module. The “STATUS_CTL” receives it's micro-coded instructions from the “STATUS_MCODE” module via the “mc_word” bus and interprets them as to the action it's to take.

Signal Name	Width	Direction	Description
rst_n	1	In	reset (active low)
clk	1	In	Clock
addr_o	8	Out	Wishbone Address
dat_i	8	In	Wishbone Data In
dat_o	8	Out	Wishbone Data Out
we_o	1	Out	Wishbone Write
sel_o	1	Out	Wishbone Select (redundant with strobe)
cyc_o	1	Out	Wishbone Cycle
ack_i	1	In	Wishbone Ack
mc_addr	3	Out	Microcode Address
mc_word	16	In	Microcode Data Word
go_ahead	1	In	Start interpreting the Microcode
stat_in	8	In	Synchronized version of masked “received_in” bus

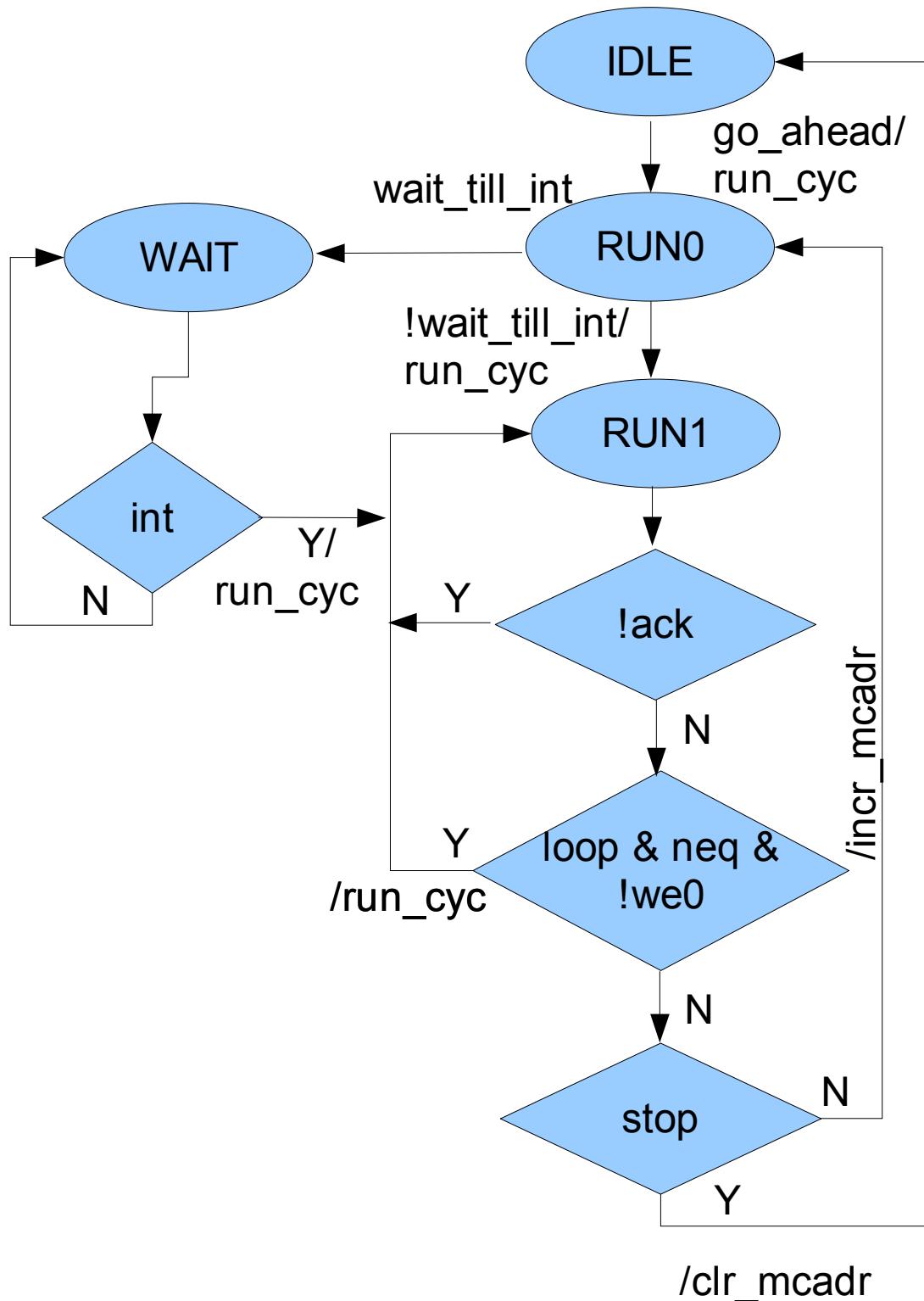
The module receives “mc_word” which controls the operation of the module. The table below defines how the state machine decodes the micro-control word.

Signal Name	Function
mc_word[7:0]	Data Out/Comparison word
mc_word[10:8]	Address to I2C Master
mc_word[11]	Source Select 0 = mc_word[7:0] 1=status_in
mc_word[14:13]	00 – Wishbone Write Operation
	01 – Wishbone Read and Compare/Loop
	10 = Wait till Interrupt
	11 – Undefined
mc_word[15]	Stop execution after transaction

This section provides a block diagram of the STAT_CTL module:



STAT_CTL Block Diagram

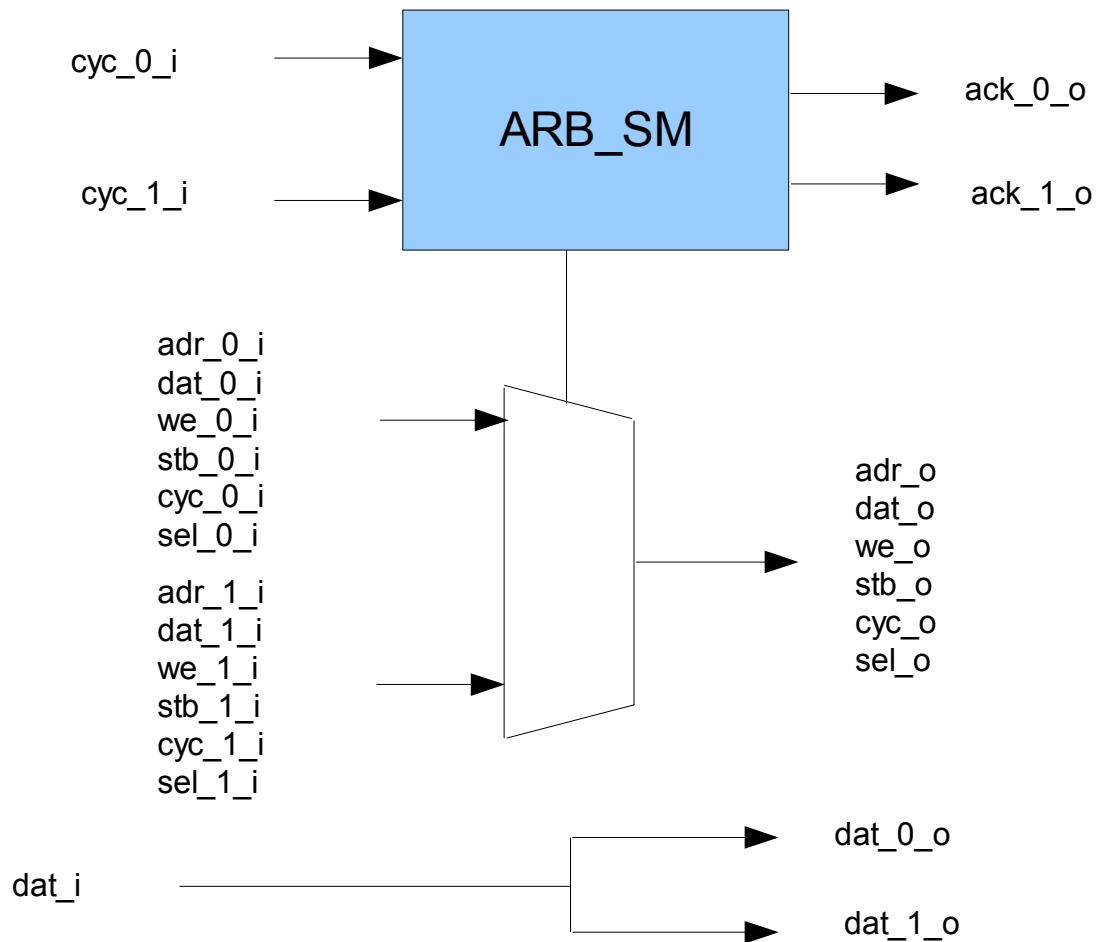


STAT_CTL State Machine

2.7 WB_ARB_MUX

This module provides arbitration services for the Wishbone masters and muxes the address bus and data_out buses.

Signal Name	Width	Direction	Description
wb_clk_i	1	In	Clock
wb_rst_i	1	In	reset (active low)
adr_0_i	8	In	Wishbone Master 0 Address
dat_0_i	8	In	Wishbone Master 0 Data In
dat_0_o	8	Out	Wishbone Master 0 Data Out
we_0_o	1	Out	Wishbone Master 0 Write
sel_0_i	1	In	Wishbone Master 0 Select (redundant with strobe)
stb_0_i	1	In	Wishbone Master 0 Strobe
ack_0_o	1	Out	Wishbone Master 0 Ack
cyc_0_i	1	In	Wishbone Master 0 Cycle
adr_1_i	8	In	Wishbone Master 1 Address
dat_1_i	8	In	Wishbone Master 1 Data In
dat_1_o	8	Out	Wishbone Master 1 Data Out
we_1_o	1	Out	Wishbone Master 1 Write
sel_1_i	1	In	Wishbone Master 1 Select (redundant with strobe)
stb_1_i	1	In	Wishbone Master 1 Strobe
ack_1_o	1	Out	Wishbone Master 1 Ack
cyc_1_i	1	In	Wishbone Master 1 Cycle
adr_o	8	Out	Wishbone Address muxed
dat_o	8	Out	Wishbone Data out (muxed)
dat_i	8	In	Wishbone Data In (muxed)
we_o	1	Out	Wishbone write (muxed)
sel_o	1	Out	Wishbone Select (muxed)
stb_o	1	Out	Wishbone Strobe (muxed)
cyc_o	1	Out	Wishbone Cycle (muxed)
ack_i	1	In	Wishbone Ack (muxed)



3 Microcode Control

This section discusses how the microcode control is used. The basic idea is that a programmed response is required, consequently an agent that can interact with I2C Master over the wishbone bus is required. The micro-word is organized to be able to control reading/writing of the bus, and has the ability to write literals to selected addresses, write the current received data bus, and to loop on a masked status.

The micro-code is organized as illustrated below:

Signal Name	Function
mc_word[7:0]	Data Out/Comparison word
mc_word[10:8]	Address to I2C Master
mc_word[11]	Source Select 0 = mc_word[7:0] 1=status_in
mc_word[14:13]	00 – Wishbone Write Operation
	01 – Wishbone Read and Compare/Loop
	10 = Wait till Interrupt
	11 – Undefined
mc_word[15]	Stop execution after transaction

The easiest way to illustrate it's use is to give an example:

This is the microcode loaded in the basic test bench. It achieves the sending of the status word to the I2C slave addressed at 0x30. Note that bit 12 in the control word isn't currently defined.

```
wr_single(SADR1, 8'h00,8'h30); // Load the Slave address into bits 7-0
wr_single(SADR1, 8'h01,{MCRUN,MCWRT,MCWORD,TXR});
```

Microword 0 = 16'b1_00_X_0_011_00110000

```
// This operation starts the first transmission of the Address phase of the I2C protocol
wr_single(SADR1, 8'h02, STA | WRB ); // Set the Start bit and Write Bit in the CTL Reg
wr_single(SADR1, 8'h03, {MCRUN,MCWRT,MCWORD,CR}); // Write to Ctl Reg -
```

Microword 1 = 16'b1_00_X_0_100_10010000

```
/ This has the affect of continually reading the status register until the ANDed condition is true,
// i.e. Bit 1 set in this case.
wr_single(SADR1, 8'h04, 8'h02); // Mask for bit 1 of the Status word
wr_single(SADR1, 8'h05, {MCRUN,MCRDC,MCWORD,SR}); // Continually read SR until mask true
```

Microword 2 = 16'b1_01_X_0_100_00000010

```
// Load the received_in bus into the I2C Master XMIT register
wr_single(SADR1, 8'h06, 8'h0); // Uses STAT_IN – so noop
wr_single(SADR1, 8'h07, {MCRUN,MCWRT,STAT_IN,TXR}); // Selects STAT_IN
```

Microword 3 = 16'b1_00_X_1_011_00000000

```
// Send the status word - and set the stop bit when done.
wr_single(SADR1, 8'h08, STO | WRB ); // Set the STOP bit and Write Bit (last operation)
wr_single(SADR1, 8'h09, {MCRUN,MCWRT,MCWORD,CR});
```

Microword 3 = 16'b1_00_X_0_100_01010000

```
// This next sequence is the last in the Microcode sequence -
// The status word is polled until bit one sets. Since MCSTOP is selected, the SM will stop
// upon completion of this step.
wr_single(SADR1, 8'h0A, 8'h02 ); // Mask for bit 1
wr_single(SADR1, 8'h0B, {MCSTOP,MCRDC,MCWORD,SR}); // Stop operation upon completion.
```

Microword 3 = 16'b0_01_X_0_100_00000010